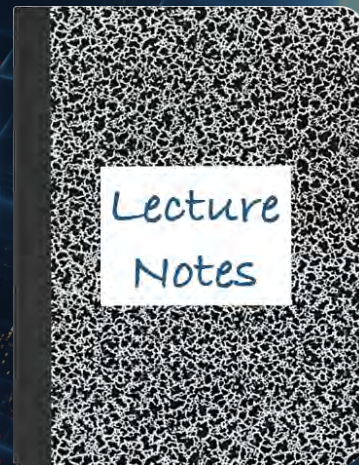


CS 417 – DISTRIBUTED SYSTEMS

Week 5: Part 4
Chubby

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Chubby

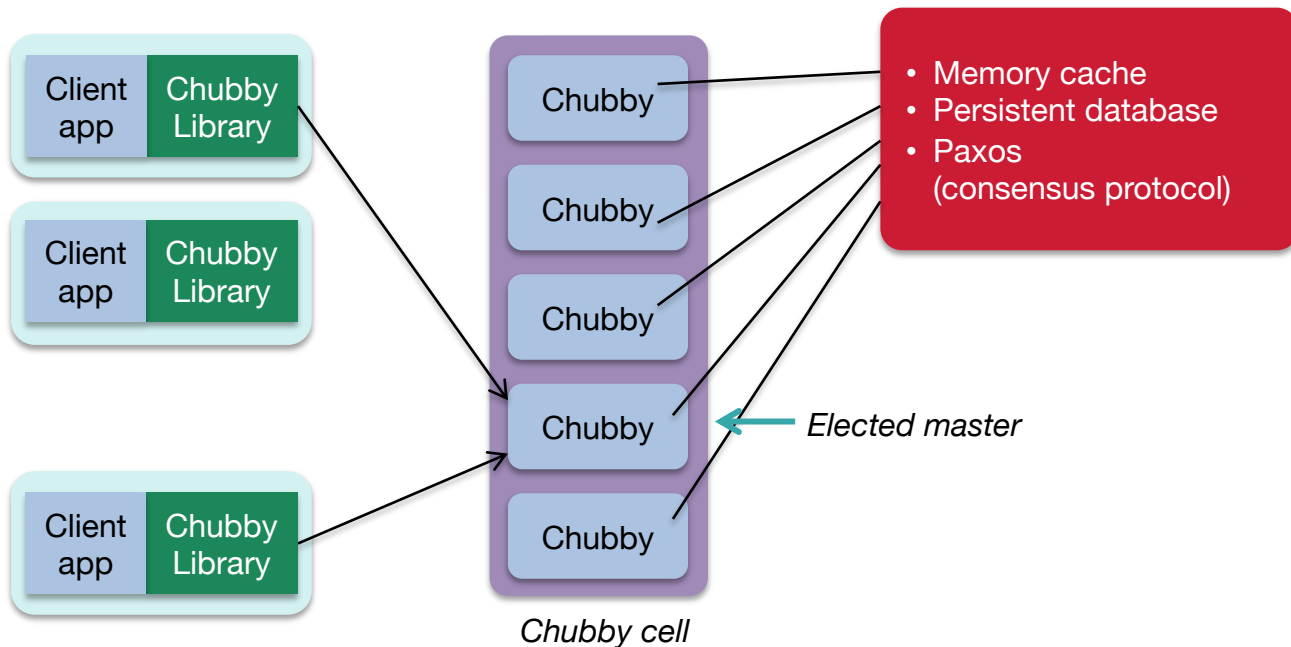
Distributed lock service + simple fault-tolerant file system

- Interfaces
 - File access
 - Event notification
 - File locking
- Chubby is used to:
 - Manage coarse-grained, **long-term locks** (hours or days, not < sec)
 - get/release/check lock – identified with a name
 - Store small **amounts of data** associated with a name
 - E.g., system configuration info, identification of primary coordinators
 - Elect masters

Design priority: availability rather than performance

Chubby Deployment

Client library + a Chubby cell (5 replica servers)



Chubby Master

- Chubby has **at most** one master
 - All requests from the client go to the master
- All other nodes (replicas) must agree on who the master is
 - Paxos consensus protocol used to elect a master
 - Master gets a lease time
 - Re-run master selection after lease time expires to extend the lease
...or if the master fails
 - When a Chubby node receives a proposal for a new master
It will accept it *only* if the old master's lease expired

Simple User-level API for Chubby

- User-level RPC interface
 - Not implemented as a true file system (e.g., under VFS in Linux)
 - Programs must access Chubby via an API
- Look up Chubby nodes via DNS
- Ask any Chubby node for the master node
- File system interface (names, content, and locks)

Chubby: File System Interface

- `/ls/cell/rest/of/name`
 - `/ls`: lock service (common to all Chubby names)
 - `cell`: resolved to a set of servers in a Chubby cell via DNS lookup
 - `/rest/of/name`: interpreted within the cell
- Each file has
 - Name
 - Data
 - Access control list
 - Lock
 - No modification, access times
 - No seek or partial reads/writes; no symbolic links; no moves

Chubby: API

<code>open()</code>	Set mode: read, write & lock, change ACL, event list, lock-delay, create
<code>close()</code>	Close access to an object
<code>GetContentsAndStat()</code>	Read file contents & metadata
<code>SetContents()</code> , <code>SetACL()</code>	Write file contents or ACL
<code>Delete()</code>	Delete an object
<code>Acquire()</code> , <code>TryAcquire()</code> , <code>Release()</code>	Lock operations
<code>GetSequencer()</code>	Sequence # for a lock
<code>SetSequencer()</code>	Associate a sequencer with a file handle
<code>CheckSequencer()</code>	Check if sequencer is valid

Chubby: Locks

- Every file & directory can act as a reader-writer lock
 - Either one client can hold an exclusive (writer) lock
 - Or multiple clients can hold reader locks
- Locks are advisory
- If a client releases a lock, the lock is immediately available
- If a client fails, the lock will be unavailable for a *lock-delay* period (typically 1 minute)

Using Locks for Leader Election

- **Using Chubby locks makes leader election easy**
 - No need for user servers to participate in a consensus protocol
... the programmer doesn't need to figure out Paxos (or Raft)
 - Chubby provides the fault tolerance
 - Participant tries to acquire a lock
 - If it gets it, then it's the master for whatever service it's providing!
- **Example: electing a master & using it to write to a file server**
 - Participant gets a lock, becomes master (*for its service, not Chubby*)
 - Gets a lock sequence count
 - In each RPC to a server, send the **sequence count** to the server
 - During request processing, a server will reject old (delayed) packets

```
if (sequence_count < current_sequence_count)
    reject request /* it must be from a delayed packet */
```

Clients may subscribe to events:

- File content modifications
- Child node added/removed/modified
- Chubby master failed over
- File handle & its lock became invalid
- Lock acquired
- Conflicting lock request from another client

Chubby client caching & master replication

- **At the client**
 - Data cached in memory by chubby clients
 - Cache is maintained by a Chubby lease, which can be invalidated
 - All clients write through to the Chubby master
- **At the master**
 - Writes are propagated via Paxos consensus to all Chubby replicas
 - Data updated in total order – replicas remain synchronized
 - The master replies to a client *after* the writes reach a **majority** of replicas
 - Reads can be acknowledged immediately without consulting the replicas.
 - Cache invalidations
 - Master keeps a list of what each client may be caching
 - Invalidations sent by master and are acknowledged by client
 - File is then cacheable again
 - Chubby database is backed up to GFS every few hours

Chubby – Apache ZooKeeper™

- Chubby is an internal Google service
- Apache ZooKeeper is an open-source centralized service for locking and storing shared services
- Built based on the Chubby paper
 - Uses Zab instead of Paxos for consensus (another protocol – roughly similar to Paxos but with a focus on log replication)
 - Replicated servers
 - Shared hierarchical namespace
 - Stored in memory and organized into files and directories
 - Locking
 - Event notification



The End